
AN EXAMPLE OF A FINITE NUMBER OF RECURSIVELY
ENUMERABLE m -DEGREES CONTAINING AN INFINITE
SEQUENCE OF RECURSIVELY ENUMERABLE SETS

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Елка Божикова. ПРИМЕР КОНЕЧНОГО ЧИСЛА РЕКУРСИВНО ПЕРЕЧИСЛЯЕМЫХ m -СТЕПЕНЕЙ, СОДЕРЖАЩИХ БЕЗКОНЕЧНУЮ ПОСЛЕДОВАТЕЛЬНОСТЬ РЕКУРСИВНО ПЕРЕЧИСЛЯЕМЫХ МНОЖЕСТВ

Цель этой статьи построить для произвольного натурального числа n рекурсивно перечисляемое (р. п.) множество A , такого чтобы множества A, A^2, \dots, A^n принадлежали разным m -степеням, а множества $A^n, A^{n+1}, A^{n+2}, \dots$ — одной и той же m -степени. Таким образом мы получаем пример для n разных р. п. m -степеней $d_m(A), d_m(A^2), \dots, d_m(A^{n-1}), d_m(A^n)$, принадлежащих к одной *btt*-степени (точнее, *c*-степени — степень $d_c(A)$ множества A совпадающая с $d_c(A^2), \dots, d_c(A^n), \dots$). Для этого достаточно доказать, что множество A^{n+1} m -сводимо к множеству A^n , а множество A^n не m -сводимо к множеству A^{n-1} . Для этой цели создаем схему для построения р. п. множества A , сопоставимого последовательности р. п. множеств A_0, A_1, A_2, \dots , так чтобы первое условие было выполнено независимо от выбора множеств A_0, A_1, A_2, \dots . После этого мы строим множества A_0, A_1, A_2, \dots по этапам, пользуясь методом приоритета, чтобы выполнить второе условие.

Elka Bojkova. AN EXAMPLE OF A FINITE NUMBER OF RECURSIVELY ENUMERABLE m -DEGREES CONTAINING AN INFINITE SEQUENCE OF RECURSIVELY ENUMERABLE SETS

The aim of this paper is to construct for an arbitrary natural number n a recursively enumerable (r. e.) set A such that the sets A, A^2, \dots, A^n belong to different m -degrees and the sets $A^n, A^{n+1}, A^{n+2}, \dots$ belong to the same m -degree. That will provide an example of n distinct r. e. m -degrees $d_m(A), d_m(A^2), \dots, d_m(A^n)$ belonging to one r. e. *btt*-degree (more precisely, *c*-degree, the degree $d_c(A)$ of the set A coincidental with $d_c(A^2), \dots, d_c(A^n), \dots$). It suffices to

prove that the set A^{n+1} is m -reducible to the set A^n , but the set A^n is not m -reducible to the set A^{n-1} . For this purpose we construct a scheme for building a r. e. set A , corresponding to a sequence of r. e. sets A_0, A_1, A_2, \dots , such that the first condition holds regardless the choice of A_0, A_1, A_2, \dots . Further we build the sets A_0, A_1, A_2, \dots by steps, using the priority argument to ensure the second condition.

In [1] Fischer proves that the reducibilities \leq_m and \leq_{btt} differ on the recursively enumerable (r.e.) non-recursive sets. He constructs an example of a r.e. btt -degree containing infinitely many distinct m -degrees. In [2] Odifreddi asks if every r.e. tt -degree contains one or infinitely many r.e. m -degrees. Downey [3] solves Odifreddi's question by constructing a r.e. tt -degree containing exactly 3 r.e. m -degrees. Rogers [4] and Odifreddi [5] summarize the obtained results about the structure of the different kinds of degrees.

In this paper we construct (for an arbitrary natural number n) a r.e. set A such that

$$A <_m A^2 <_m \dots <_m A^n \equiv_m A^{n+1} \equiv_m A^{n+2} \equiv_m \dots$$

That will provide an example of n distinct r.e. m -degrees ($d_m(A), d_m(A^2), \dots, d_m(A^{n-1})$ and $d_m(A^n)$) belonging to one r.e. btt -degree (more precisely, c -degree — the degree $d_c(A)$ of the set A coincidental with $d_c(A^2), \dots, d_c(A^n), \dots$).

Obviously, all the recursive sets belong to one tt -degree and one m -degree, i.e. for any recursive set B the following equivalencies hold:

$$B \equiv_m B^2 \equiv_m B^3 \equiv_m \dots$$

Now, we have to find a set A such that the sets A, A^2, \dots, A^n belong to different m -degrees and the sets $A^n, A^{n+1}, A^{n+2}, \dots$ belong to the same m -degree. It suffices for this purpose to prove

$$(1) \quad A^{n+1} \leq_m A^n$$

and

$$(2) \quad \neg A^n \leq_m A^{n-1}.$$

The method elaborated by Ditchov in [6] is used to construct the set responding to both conditions. First, we construct a scheme for building a r.e. set A corresponding to a sequence of r.e. sets A_0, A_1, A_2, \dots (answering to some conditions), such that (1) holds regardless the choice of A_0, A_1, A_2, \dots . After that we build the sets A_0, A_1, A_2, \dots by steps, using the priority argument to ensure (2).

* * *

(1.1) **Definition.** Let A and B be sets of natural numbers:

a) We say that the set A is " m -reducible" to the set B ($A \leq_m B$) if there exists a recursive function $f : N \rightarrow N$ such that

$$\forall x (x \in A \iff f(x) \in B);$$

b) We say that the sets A and B are “ m -equivalent” ($A \equiv_m B$) if

$$A \leq_m B \quad \text{and} \quad B \leq_m A.$$

We shall denote by “ $A <_m B$ ” the fact that $A \leq_m B$ but $A \not\equiv_m B$.

Let for any natural number n the recursive functions $J^n, J_1^n, J_2^n, \dots, J_n^n$ be such that for any n and for any x there exists a single tuple x_1, \dots, x_n such that

$$(1.2) \quad x = J^n(x_1, \dots, x_n)$$

and

$$(1.3) \quad \forall i_{1 \leq i \leq n} [J_i^n(J^n(x_1, \dots, x_n)) = x_i].$$

(1.4) **Definition.** Let A be an arbitrary set of natural numbers. For any natural number n we define the set A^n :

$$A^n = \{J^n(x_1, \dots, x_n) \mid x_1 \in A \& \dots \& x_n \in A\}.$$

Obviously, for any set A

$$A \leq_m A^2 \leq_m \dots \leq_m A^n \leq_m A^{n+1} \leq_m A^{n+2} \leq_m \dots,$$

i. e. the essential part of the problem is to find a set A such that:

(1) for every $i, 1 \leq i \leq n-1, \neg A^{i+1} \leq_m A^i$;

(2) for every $i, i \geq n, A^{i+1} \leq_m A^i$.

It is easy to prove the following two lemmas:

(1.5) **Lemma.** Let A be an arbitrary subset of N and n be an arbitrary natural number. Then $A^{n+1} \leq_m A^n$ iff there exist recursive functions $f_1 : N^{n+1} \rightarrow N, f_2 : N^{n+1} \rightarrow N, \dots, f_n : N^{n+1} \rightarrow N$ such that

$$(1.6) \quad x_1 \in A \& \dots \& x_{n+1} \in A$$

$$\iff f_1(x_1, \dots, x_{n+1}) \in A \& \dots \& f_n(x_1, \dots, x_{n+1}) \in A.$$

Proof. Using the fact that one unary function g reduces A^{n+1} to A^n iff for the $(n+1)$ -ary functions f_1, \dots, f_n , defined as follows:

$$f_i(x_1, x_2, \dots, x_{n+1}) = J_i^n(g(J^{n+1}(x_1, x_2, \dots, x_{n+1}))), \quad i = 1, 2, \dots, n,$$

the condition (1.6) holds, one can easily verify that Lemma (1.5) is true.

(1.7) **Lemma.** a) If $A^{n-1} <_m A^n$, then for any $m, 2 \leq m \leq n-1$:

$$A^{m-1} <_m A^m;$$

b) If $A^n \equiv_m A^{n+1}$, then for any $m, m > n$:

$$A^m \equiv_m A^{m+1}.$$

Hence to solve our problem, it is enough to prove

$$A^{n-1} <_m A^n \equiv_m A^{n+1}.$$

Let n be an arbitrary natural number*, $n > 1$. We shall find a set A such that

$$A <_m A^2 <_m \dots <_m A^n \equiv_m A^{n+1} \equiv_m A^{n+2} \equiv_m \dots$$

(2.1) **Definition.** Let L_n be a language with the following alphabet:

- c_0, c_1, \dots — infinite sequence of constants;
- x_1, \dots, x_n — n variables;
- F_1, \dots, F_n — $(n + 1)$ -ary functional symbols.

Terms are defined by means of the following inductive clauses:

- a) c_i is a term, $i \in N$; x_i is a term, $1 \leq i \leq n$;
- b) if $\tau^1, \dots, \tau^{n+1}$ are terms, then $F_i(\tau^1, \dots, \tau^{n+1})$ is a term, $1 \leq i \leq n$.

(2.2) **Definition.** For any term τ we define *subterms* of τ :

- a) if $\tau = c_i$, $i \in N$, c_i is a subterm of τ ; if $\tau = x_i$, $1 \leq i \leq n$, x_i is a subterm of τ ;
- b) if $\tau = F_i(\tau^1, \dots, \tau^{n+1})$, $1 \leq i \leq n$, then the term τ and the subterms of $\tau^1, \dots, \tau^{n+1}$ are subterms of τ .

(2.3) **Definition.** For any term τ we define *deepness* $dp(\tau)$ of τ :

- a) if $\tau = c_i$, $i \in N$ or $\tau = x_i$, $1 \leq i \leq n$, then $dp(\tau) = 1$;
- b) if $\tau = F_i(\tau^1, \dots, \tau^{n+1})$, $1 \leq i \leq n$, and $dp(\tau^1), \dots, dp(\tau^{n+1})$ are defined, then $dp(\tau) = 1 + \max_{1 \leq j \leq n+1} dp(\tau^j)$.

(2.4) **Definition.** We will call a *partial structure* every ordered $(n + 1)$ -tuple $U = \langle N; \theta_1, \dots, \theta_n \rangle$, where $\theta_1, \dots, \theta_n$ are partial functions of $n + 1$ variables.

(2.5) **Definition.** Let $U = \langle N; \theta_1, \dots, \theta_n \rangle$ be a partial structure. We define the *value* τ_U of the term τ in the partial structure U :

- a) if $\tau = c_i$, $\tau_U(t_1, \dots, t_n) = i$ for any t_1, \dots, t_n , $i \in N$;
if $\tau = x_i$, $\tau_U(t_1, \dots, t_n) = t_i$ for any t_1, \dots, t_n , $1 \leq i \leq n$;
- b) if $\tau = F_i(\tau^1, \dots, \tau^{n+1})$, $1 \leq i \leq n$, $t_1 \in N, \dots, t_n \in N$ and $\tau_U^1, \dots, \tau_U^{n+1}$ are defined, then

$$\tau_U(t_1, \dots, t_n) \cong \theta_i(\tau_U^1(t_1, \dots, t_n), \dots, \tau_U^{n+1}(t_1, \dots, t_n)).$$

(2.6) **Definition.** We define the *cod* $cd(\tau)$ of the term τ :

- a) if $\tau = c_i$, $i \in N$, $cd(\tau) = J^2(0, i)$; if $\tau = x_i$, $1 \leq i \leq n$, $cd(\tau) = J^2(1, i)$;
- b) if $\tau = F_i(\tau^1, \dots, \tau^{n+1})$, $1 \leq i \leq n$, $cd(\tau) = J^{n+2}(i + 1, cd(\tau^1), \dots, cd(\tau^{n+1}))$.

That coding allows us to verify for any natural number whether it is a cod of any term and if so, to find this term.

Theorem 1. Let $M = \{2k \mid k \in N\}$. There exists a partial structure $U = \langle N, \theta_1, \dots, \theta_n \rangle$ such that the functions $\theta_1, \dots, \theta_n$ are recursive and:

* n is fixed till the end of this paper.

$$1) \forall t_1 \in N \dots \forall t_{n+1} \in N,$$

$$(*) \quad (t_1 \in M \& \dots \& t_{n+1} \in M \\ \iff \theta_1(t_1, \dots, t_{n+1}) \in M \& \dots \& \theta_n(t_1, \dots, t_{n+1}) \in M);$$

$$2) \forall (\tau^1, \dots, \tau^{n-1} - \text{terms}) \exists t_1 \in N, \dots, t_n \in N,$$

$$(**) \quad \neg (t_1 \in M \& \dots \& t_n \in M \\ \iff \tau_U^1(t_1, \dots, t_n) \in M \& \dots \& \tau_U^{n-1}(t_1, \dots, t_n) \in M)).$$

Proof. We shall say that for the partial functions $\omega_1, \dots, \omega_n$ the condition $(\tilde{*})$ holds if $\text{dom } \omega_1 = \dots = \text{dom } \omega_n$ and the condition $(*)$ holds for $\text{dom } \omega_1$.

We build the functions $\theta_1, \dots, \theta_n$ by steps: at any step $s = 0, 1, 2, \dots$ we build the functions $\theta_1^{(s)}, \dots, \theta_n^{(s)}$ as finite extensions of $\theta_1^{(s-1)}, \dots, \theta_n^{(s-1)}$ such that for $\theta_1^{(s)}, \dots, \theta_n^{(s)}$ the condition $(\tilde{*})$ holds. For $s = 0$ we accept $\theta_1^{(-1)} = \dots = \theta_n^{(-1)} = \emptyset$. (The domain of \emptyset is empty.)

$$\text{Finally, } \theta_i = \bigcup_{s \in N} \theta_i^{(s)}, \quad i = 1, \dots, n.$$

I. By the even steps $s = 2p$ we shall ensure that the functions $\theta_1^{(s)}, \dots, \theta_n^{(s)}$ are defined in $(J_1^{n+1}(p), \dots, J_{n+1}^{n+1}(p))$ and therefore the functions $\theta_1, \dots, \theta_n$ will be total.

1) If $(J_1^{n+1}(p), \dots, J_{n+1}^{n+1}(p)) \notin \text{dom } \theta_1^{(s-1)}$, then for all $l, 1 \leq l \leq n$,

$$\theta_l^{(s)}(J_1^{n+1}(p), \dots, J_{n+1}^{n+1}(p)) = \begin{cases} 2, & \text{if } J_1^{n+1}(p) \in M \& \dots \& J_{n+1}^{n+1}(p) \in M, \\ 3, & \text{otherwise,} \end{cases}$$

and

$$\theta_l^{(s)}(y_1, \dots, y_{n+1}) \cong \theta_l^{(s-1)}(y_1, \dots, y_{n+1}), \quad \text{if } J^{n+1}(y_1, \dots, y_{n+1}) \neq p.$$

2) If $(J_1^{n+1}(p), \dots, J_{n+1}^{n+1}(p)) \in \text{dom } \theta_1^{(s-1)}$, then $\forall l_{1 \leq l \leq n} [\theta_l^{(s)} = \theta_l^{(s-1)}]$.

II. By the odd steps $s = 2p + 1, p \in N$, if there are not terms $\tau^1, \dots, \tau^{n-1}$ such that $p = J^{n-1}(\text{cd}(\tau^1) \dots, \text{cd}(\tau^{n-1}))$, we take $\forall l_{1 \leq l \leq n} [\theta_l^{(s)} = \theta_l^{(s-1)}]$. If such terms exist, we shall find t_1, \dots, t_n such that the condition $(**)$ holds. In addition we obtain that the condition $(**)$ holds for any terms $\tau^1, \dots, \tau^{n-1}$.

In the second case we first verify whether some of the terms $\tau^1, \dots, \tau^{n-1}$ is a constant with a value — an odd number. If so, for any n -tuple of even numbers t_1, \dots, t_n the condition $(**)$ holds and we do nothing; i. e. we take $\theta_l^{(s)} = \theta_l^{(s-1)}$ for all $l, 1 \leq l \leq n$. If all the terms $\tau^1, \dots, \tau^{n-1}$, which are constants, have the values — even numbers, we shall ensure $(**)$ by finding t_1, \dots, t_n and building $\theta_1^{(s)}, \dots, \theta_n^{(s)}$ such that:

$$\begin{aligned} & - \tau_{U^{(s)}}^1(t_1, \dots, t_n) \in M \& \dots \& \tau_{U^{(s)}}^{n-1}(t_1, \dots, t_n) \in M; \\ & - t_1 \notin M \vee \dots \vee t_n \notin M. \end{aligned}$$

Every term is a constant, a variable, or its deepness is bigger than 1.

(1) For all the terms $\tau^1, \dots, \tau^{n-1}$, which are constants, we have that their values belong to M .

(2) For these ones, which are variables, we shall choose the corresponding numbers t_1, \dots, t_n (which are values of such terms) to be even and these terms also will satisfy the condition. Because there is at least one number among t_1, \dots, t_n which is not a value of any term, let this number be odd and by this way t_1, \dots, t_n also will satisfy the condition. In both cases we shall take t_1, \dots, t_n big enough not to belong to the domains of $\theta_1^{(s-1)}, \dots, \theta_n^{(s-1)}$ and thus to the values of the terms with deepness bigger than 1.

(3) For the terms with deepness bigger than 1 we need an auxiliary lemma — Lemma 1. It will be applied for the partial structure

$$U^{(s-1)} = \langle N, \theta_1^{(s-1)}, \dots, \theta_n^{(s-1)} \rangle \text{ for } \tau^1, \dots, \tau^{n-1}, t_1, \dots, t_n$$

(obtained from (2) or, if there are no terms with deepness 1, we take t_1, \dots, t_n big enough and such that $t_1 \notin M \vee \dots \vee t_n \notin M$) and for one term τ among the terms $\tau^1, \dots, \tau^{n-1}$, $\text{dp}(\tau) > 1$. Using this lemma we shall obtain the functions $\theta_1^{(s)}, \dots, \theta_n^{(s)}$ such that $\theta_i^{(s-1)} \leq \theta_i^{(s)}$, $1 \leq i \leq n$, the condition $(\tilde{*})$ holds and the value of the term τ for (t_1, \dots, t_n) is an even number. If at the same time (t_1, \dots, t_n) enters in the domain of the value of some other term among $\tau^1, \dots, \tau^{n-1}$ (except τ), we have to ensure that the value of the other term in (t_1, \dots, t_n) is also an even number.

Lemma 1. *Let the terms $\tau^1, \dots, \tau^{n-1}$ be given. For any partial structure $U = \langle N; \theta_1, \dots, \theta_n \rangle$ with finite functions satisfying $(\tilde{*})$, for any term τ with $\text{dp}(\tau) \geq 2$, for any natural number m , and for any n -tuple of natural numbers (t_1, \dots, t_n) such that*

$$(t_1, \dots, t_n) \notin \text{dom } \tau_U \text{ and } t_1 \notin M \vee \dots \vee t_{n+1} \notin M$$

there exists a partial structure $U' = \langle N, \theta'_1, \dots, \theta'_n \rangle$ with finite recursive functions satisfying $(\tilde{})$, such that:*

- (i) $\theta'_i \geq \theta_i$ for every i , $1 \leq i \leq n$;
- (ii) $(t_1, \dots, t_n) \in \text{dom } \tau_{U'}$ and $\tau_{U'}(t_1, \dots, t_n) > z$,

where $z = \max(\{m, \max(I_1(\text{dom } \theta'_1)), \dots, \max(I_{n+1}(\text{dom } \theta'_1))\})$;

- (iii) if $(t_1, \dots, t_n) \in (\text{dom } \tau_{U'}^j \setminus \text{dom } \tau_U^j)$, then $\tau_{U'}^j(t_1, \dots, t_n) \in M$, $j = 1, \dots, n-1$.

We shall apply this lemma successively for all the terms among $\tau^1, \dots, \tau^{n-1}$ with deepness bigger than 1 and we shall obtain

$$\tau_{U^{(s)}}^1(t_1, \dots, t_n) \in M \ \& \ \dots \ \& \ \tau_{U^{(s)}}^{n-1}(t_1, \dots, t_n) \in M;$$

$$t_1 \notin M \vee \dots \vee t_n \notin M.$$

The idea to prove Lemma 1 is to find for the term $\tau = F_i(\sigma^1, \dots, \sigma^{n+1})$ a partial structure $V = \langle N, \theta''_1, \dots, \theta''_{n+1} \rangle$ such that $\forall j$, $1 \leq j \leq n+1$, $(t_1, \dots, t_n) \in \text{dom } \sigma_V^j$, $z_j = \sigma_V^j(t_1, \dots, t_n)$ and $(z_1, \dots, z_{n+1}) \notin \text{dom } \theta''_{n+1}$. In this case we shall build the functions $\theta'_1, \dots, \theta'_{n+1}$ such that $\theta'_1 \geq \theta''_1, \dots, \theta'_{n+1} \geq \theta''_{n+1}$, $\text{dom } \theta'_{n+1} =$

$\text{dom } \theta''_{n+1} \cup \{(z_1, \dots, z_{n+1})\}$, and the values of the functions in (z_1, \dots, z_{n+1}) are defined as follows (z is big enough):

1) If $z_1 \in M \& \dots \& z_{n+1} \in M$, the values of all the functions are $2z$;

2) If $z_1 \notin M \vee \dots \vee z_{n+1} \notin M$, and for some term τ^j among $\tau^1, \dots, \tau^{n-1}$ there exist a number l , $1 \leq l \leq n$, and terms $\varepsilon^1, \dots, \varepsilon^{n+1}$ such that $\tau^j = F_l(\varepsilon^1, \dots, \varepsilon^{n+1})$, then the value of the function θ'_l is also $2z$;

3) In the other cases the values of the functions are $2z + 1$.

So we ensure that $(\bar{*})$ is true for the new partial structure (if $z_1 \notin M \vee \dots \vee z_{n+1} \notin M$, because the terms $\tau^1, \dots, \tau^{n-1}$ are $n - 1$ and the symbols F_1, \dots, F_n are n , at least one of the functions is in the case 3) and its value is $2z + 1$). Except that we ensure $\tau_{U'}(t_1, \dots, t_n)$ to be big enough and for any of the terms τ^j , $j = 1, \dots, n-1$, such that $(t_1, \dots, t_n) \in (\text{dom } \tau_{U'}^j \setminus \text{dom } \tau_U^j)$, $\tau_{U'}^j(t_1, \dots, t_n) \in M$ is true. The last follows from 2), but not evidently. So we need an auxiliary lemma — Lemma (2.7).

With this lemma we have to prove that for some l , $1 \leq l \leq n$, $\tau_{U'}^j(t_1, \dots, t_n) = \theta_1(z_1, \dots, z_{n+1})$ and θ_1 is in case 2), and therefore $\tau_{U'}^j(t_1, \dots, t_n) = 2z \in M$. In the general case the lemma can not be proved for the term τ^j , it is only true that there exists a subterm of τ^j for which the condition holds, but in this concrete application of Lemma (2.7) in Lemma 1 we can prove that this subterm may be only the term τ^j .

(2.7) Lemma. *For any two partial structures $U = \langle N, \theta_1, \dots, \theta_n \rangle$ and $U' = \langle N, \theta'_1, \dots, \theta'_n \rangle$, for any term σ and for any natural numbers y_1, \dots, y_{n+1} ; t_1, \dots, t_n , such that:*

a) $\text{dom } \theta_1 \equiv \dots \equiv \text{dom } \theta_n$ and that is a finite set;

b) $\text{dom } \theta'_i = \text{dom } \theta_i \cup \{(y_1, \dots, y_{n+1})\}$, $i = 1, \dots, n$, and $(y_1, \dots, y_{n+1}) \notin \text{dom } \theta_1$;

c) $\theta_1 \leq \theta'_1, \dots, \theta_n \leq \theta'_n$;

d) $(t_1, \dots, t_n) \in (\text{dom } \sigma_{U'} \setminus \text{dom } \sigma_U)$,

there exist a subterm σ' of σ , $i \in \{1, \dots, n\}$, and terms $\varepsilon^1, \dots, \varepsilon^{n+1}$ such that:

$$(i) \quad \sigma' = F_i(\varepsilon^1, \dots, \varepsilon^{n+1});$$

$$(ii) \quad \varepsilon_{U'}^1(t_1, \dots, t_n) = y_1, \dots, \varepsilon_{U'}^{n+1}(t_1, \dots, t_n) = y_{n+1}.$$

Proof. By the definitions (2.3) and (2.5) we have

(2.8) $\forall \tau$ -term, $\forall U$ -partial structure $[\text{dp}(\tau) = 1 \Rightarrow \tau_U$ is a total function].

From condition d) we have $(t_1, \dots, t_n) \notin \text{dom } \sigma_U$. Therefore $\text{dp}(\tau) \geq 2$ and from (2.3) it follows that there exist k , $1 \leq k \leq n$, and terms $\sigma^1, \dots, \sigma^{n+1}$ such that $F_k(\sigma^1, \dots, \sigma^{n+1}) = \sigma$.

We have two cases:

C a s e I. $\forall j_{1 \leq j \leq n+1} [(t_1, \dots, t_n) \in \text{dom } \sigma_U^j]$. Let $z_j = \sigma_U^j(t_1, \dots, t_n)$, $j = 1, \dots, n+1$. From d) it follows that $(z_1, \dots, z_{n+1}) \notin \text{dom } \theta_k$ and $(z_1, \dots, z_{n+1}) \in \text{dom } \theta'_k$ and from b) we obtain $z_1 = y_1, \dots, z_{n+1} = y_{n+1}$. From c) it follows that

$$\forall j_{1 \leq j \leq n+1} [\sigma_{U'}^j(t_1, \dots, t_n) = \sigma_U^j(t_1, \dots, t_n)].$$

We obtained $\sigma_{U'}^1(t_1, \dots, t_n) = y_1, \dots, \sigma_{U'}^{n+1}(t_1, \dots, t_n) = y_{n+1}$, i. e. the conditions (i) and (ii) hold for $\sigma' = \sigma$ and $\varepsilon^j = \sigma^j$, $j = 1, \dots, n+1$.

C a s e II. We apply induction on $\text{dp}(\sigma)$:

1) $\text{dp}(\sigma) = 2$. Then $\text{dp}(\sigma^1) = \dots = \text{dp}(\sigma^{n+1}) = 1$ and the lemma follows from (2.8) and Case I;

2) Let it be true for the terms with deepness less than $\text{dp}(\sigma)$;

3) We shall prove it for $\sigma = F_k(\sigma^1, \dots, \sigma^{n+1})$. The terms $\sigma^1, \dots, \sigma^{n+1}$ have smaller deepness. We have $(t_1, \dots, t_n) \notin \text{dom } \sigma_U$. Two cases are possible:

Case 1. $\exists j_{1 \leq j \leq n+1} [(t_1, \dots, t_n) \notin \text{dom } \sigma_U^j]$. But $(t_1, \dots, t_n) \in \text{dom } \sigma_{U'} \Rightarrow (t_1, \dots, t_n) \in \text{dom } \sigma_{U'}^j$, and by the induction hypothesis for σ^j we obtain that there exists a σ' -subterm of σ^j , $i, \varepsilon^1, \dots, \varepsilon^{n+1}$ such that $\sigma' = F_i(\varepsilon^1, \dots, \varepsilon^{n+1})$ and $\forall m_{1 \leq m \leq n+1} [\varepsilon_{U'}^m(t_1, \dots, t_n) = y_m]$. But σ' is a subterm of σ^j , so σ' is a subterm of σ also and in this case the lemma is proved.

Case 2. $\forall j_{1 \leq j \leq n+1} [(t_1, \dots, t_n) \in \text{dom } \sigma_U^j]$. Then Lemma (2.7) follows from Case I.

* * *

(3.1) **Definition.** The total functions $\varphi_1, \dots, \varphi_n$ of $n+1$ variables are defined as follows:

$$\varphi_i(t_1, \dots, t_{n+1}) = J^{n+2}(i, t_1, \dots, t_{n+1}), \quad i = 1, \dots, n; \quad t_1 \in N, \dots, t_{n+1} \in N.$$

(3.2) **Definition.** Let $N_0 = N \setminus (\bigcup_{i=1}^n \text{Range}(\varphi_i))$.

(3.3) **Definition.** Let $\{A_i\}_{i \in N}$ be a sequence of disjoint subsets of N_0 . The sequence $\{[A_i]\}_{i \in N}$ of disjoint subsets of N and the set A are defined as follows:

- a) if $p \in A_i$, then $p \in [A_i]$, $p \in N_0$, $i \in N$;
- b) if $p_1 \in [A_{i_1}] \& \dots \& p_{n+1} \in [A_{i_{n+1}}] \& \theta_k(i_1, \dots, i_{n+1}) = m$, $1 \leq k \leq n$, then $\varphi_k(p_1, \dots, p_{n+1}) \in [A_m]$;
- c) $A = \bigcup_{i \in M} [A_i]$. (We remind that $M = \{2k \mid k \in N\}$.)

Note. If the set $\{(p, i) \mid p \in A_i\}$ is r.e., then A is also r.e. In this case we say that we have a r.e. sequence of r.e. sets.

We can prove the following lemma:

(3.4) **Lemma.** Let $\{A_i\}_{i \in N}$ be a sequence of disjoint subsets of N_0 and the set A be obtained by Definition (3.3). Then for any $n+1$ natural numbers p_1, \dots, p_{n+1} it is true that

$$p_1 \in A \& \dots \& p_{n+1} \in A \iff \varphi_1(p_1, \dots, p_{n+1}) \in A \& \dots \& \varphi_n(p_1, \dots, p_{n+1}) \in A.$$

(3.5) **Corollary.** Let $\{A_i\}_{i \in N}$ be a sequence of disjoint subsets of N_0 and the set A be obtained by Definition (3.3). Then $A^{n+1} \leq_m A^n$.

(3.6) **Definition.** We define a correspondence between two terms τ and σ as follows:

- a) if τ has not a subterm that is a constant, then τ corresponds with $\sigma = \tau$;
- b) if c_{i_1}, \dots, c_{i_k} are all the subterms of τ which are constants, $i_1 \in [A_{r_1}], \dots, i_k \in [A_{r_k}]$, then τ corresponds with the term σ , where σ is obtained from τ by replacing c_{i_j} with c_{r_j} for any $j \in \{1, \dots, k\}$.

We need this correspondence to have the following

(3.7) **Lemma.** If $\{A_i\}_{i \in N}$ is a sequence of disjoint subsets of N_0 , τ is a term, p_1, \dots, p_n are arbitrary natural numbers, and:

a) for any constant c_j which is a subterm of τ we have a number $i \in N$ such that $j \in [A_i]$;

b) τ corresponds with σ by Definition (3.6);

and

c) $p_1 \in [A_{i_1}] \& \dots \& p_n \in [A_{i_n}] \& \sigma_U(i_1, \dots, i_n) = m$,

then $\tau_V(p_1, \dots, p_n) \in [A_m]$.

(3.8) **Lemma.** Let $V = \langle N, \varphi_1, \dots, \varphi_n \rangle$. For any natural number x there is an effective way to find a term τ that have not subterms which are variables (for $x \notin N_0$, $\tau \neq \text{constant } c_x$) such that $x = \tau_V(x_1, \dots, x_n)$ and the values of all the constants — subterms of τ , belong to N_0 .

We prove this lemma by defining the function $\|z\|$ for any $z \in N$:

- 1) if $z \in N_0$, then $\|z\| = 0$;
- (3.9) 2) if $z = \varphi_1(z_1, \dots, z_{n+1})$, then $\|z\| = 1 + \max_{1 \leq j \leq n+1} \|z_j\|$, $1 \leq i \leq n$,

and applying induction on $\|x\|$.

We shall build a r.e. sequence of disjoint r.e. subsets A_0, A_1, A_2, \dots of N_0 such that for the set A obtained by Definition (3.3) it holds

$$(3.10) \quad A^n \not\equiv_m A^{n-1}.$$

Then it follows from Lemma (1.7) and Corollary (3.5) that A is the set we need.

We shall build the sets $\{A_i\}_{i \in N}$ by steps — at any step s we build $\{A_i^{(s)}\}_{i \in N}$, ensuring that $(A^{(s)})^n$ is not m -reducible to $(A^{(s)})^{n-1}$ by the e -th recursive function, $e = J_1^2(s)$. ($A^{(s)}$ is obtained from $\{A_i^{(s)}\}_{i \in N}$ by Definition (3.3).)

At the end we take $A_i = \bigcup_{s=0}^{\infty} A_i^{(s)}$, $i \in N$. For this purpose at any step s we shall find numbers x_1, \dots, x_n such that if φ_e is the e -th partial recursive function, $e = J_1^2(s)$, $J^{(n)}(x_1, \dots, x_n) \in \text{dom } \varphi_e$ and $\varphi_e(J^{(n)}(x_1, \dots, x_n)) = J^{(n-1)}(z_1, \dots, z_{n-1})$, one of both conditions holds:

- (i) $x_1 \in A \& \dots \& x_n \in A \& \exists i_{1 \leq i \leq n-1} (z_i \notin A)$;
- (ii) $(x_1 \notin A \vee \dots \vee x_n \notin A) \& z_1 \in A \& \dots \& z_{n-1} \in A$.

If for this purpose we put the numbers x_1, \dots, x_n in some sets A_{i_1}, \dots, A_{i_n} , we create a positive e -requirement $\{x_1, \dots, x_n\}$, and if some numbers y_1, \dots, y_k

must not belong to some set, we create a negative e -requirement $\{y_1, \dots, y_k\}$. We shall use the priority argument: if at the step s we need one number x to belong to some set and at a step t — not to belong, the smaller between $J_1^2(s)$ and $J_1^2(t)$ has a priority. So, when we choose x_1, \dots, x_n at the step s , they must not belong to any negative requirements created at some steps $t < s$ such that $J_1^2(t) < J_1^2(s)$, but they may belong to negative requirements created at steps $r < s$ such that $J_1^2(r) > J_1^2(s)$. In the second case the $J_1^2(r)$ -requirement is injured and we need at some later step r' , $J_1^2(r') = J_1^2(r)$, to create a new $J_1^2(r)$ -requirement.

If one $J_1^2(s)$ -requirement is not injured at a step $r > s$, it is called active at this step. If it is active at every step $r > s$, it is called constant. At the end we shall prove that for any e the condition ((i) \vee (ii)) is injured only finite times.

Now we shall describe the construction of the sets $\{A_i\}_{i \in N}$.

(3.11) S t e p $s = 0$. Let $N_0 = N_1 \cup N_2$, where N_1 and N_2 are infinite disjoint recursive sets and $N_2 = \{a_0 < a_1 < \dots\}$. Let r' be a monotonically increasing function such that $\text{Ran}(r') = N_1$ and $r(x) = r'(n \cdot 2^x + x)$. Let

$$\varphi_{e,s}(x) \cong \begin{cases} \varphi_e(x), & \text{if } x \in \text{dom } \varphi_e \text{ and } \varphi_e(x) \text{ is countable for less than } s \text{ steps,} \\ \text{not defined,} & \text{otherwise.} \end{cases}$$

Let $A_i^{(0)} = \{a_i\}$, $i \in N$. So, all the sets are not empty.

S t e p $s > 0$. $e = J_1^2(s)$. First we verify whether there exists an active e -requirement. If such a requirement exists, then we do nothing, i.e. we take $A_i^{(s)} = A_i^{(s-1)}$, $i \in N$. Otherwise we verify whether there exist $x_1 \in N_1, \dots, x_n \in N_1, x_1 > r(e) \& \dots \& x_n > r(e)$, $J^n(x_1, \dots, x_n) \in \text{dom } \varphi_{e,s}$, belonging neither to $\bigcup_{i \in N} A_i^{(s-1)}$ nor to any active negative requirement, created at a step $t < s$ such that $J_1^2(t) < J_1^2(s)$. If such numbers do not exist, we do nothing. If there are such numbers x_1, \dots, x_n , we take the smallest — $x_1^{(e)}, \dots, x_n^{(e)}$. From the choice of J^n there exist z_1, \dots, z_{n-1} such that

$$\varphi_e(J^n(x_1^{(e)}, \dots, x_n^{(e)})) = J^{n-1}(z_1, \dots, z_{n-1}).$$

It follows from Lemma (3.8), applied for z_1, \dots, z_{n-1} , that there exist terms $\psi^1, \dots, \psi^{n-1}$ such that

$$(3.12) \quad z_i = \psi_V^i(x_1, \dots, x_n), \quad i = 1, \dots, n-1.$$

We consider these numbers among z_1, \dots, z_{n-1} for which the constants-subterms of the corresponding terms $\psi^1, \dots, \psi^{n-1}$ already belong to some sets $A_i^{(s-1)}$, $i \in N$. Let that be z_1, \dots, z_q .

C a s e I. $\exists i_{1 \leq i \leq q} \exists k \in N (z_i \in [A_{2k+1}^{(s-1)}])$, i.e. $z_i \notin A$ and in this case we satisfy (i). Let $A_{2j}^{(s)} = A_{2j}^{(s-1)} \cup \{x_j^{(e)}\}$, $j = 1, \dots, n$, and $A_l^{(s)} = A_l^{(s-1)}$ for $l \notin \{2, 4, 6, \dots, 2n\}$. We create a positive e -requirement $\{x_1^{(e)}, \dots, x_n^{(e)}\}$, which is also constant.

C a s e II. $z_1 \in A^{(s)} \& \dots \& z_q \in A^{(s)}$.

II.1. $q = n - 1$. Then we satisfy (ii).

Let $A_j^{(s)} = A_j^{(s-1)} \cup \{x_j^{(e)}\}$, $j = 1, \dots, n$, and $A_l^{(s)} = A_l^{(s-1)}$ for $l \notin \{1, \dots, n\}$.

We create a positive e -requirement $\{x_1^{(e)}, \dots, x_n^{(e)}\}$, which is also constant.

II.2. $q < n - 1$.

II.2.1. $\exists p_{1 \leq p \leq n-1}$ (ψ_p has a subterm-constant with a value y_j^p ,

$$y_j^p \notin \left(\bigcup_{i \in N} A_i^{(s-1)} \cup \{x_1^{(e)}, \dots, x_n^{(e)}\} \right)).$$

For simplicity let $p = n - 1$. We satisfy (i) — we put $x_1^{(e)}, \dots, x_n^{(e)}$ into A and create a negative requirement such that z_{n-1} stays always out of A .

We take $A_{2j}^{(s)} = A_{2j}^{(s-1)} \cup \{x_j^{(e)}\}$, $j = 1, \dots, n$, and $A_l^{(s)} = A_l^{(s-1)}$ for $l \notin \{2, 4, 6, \dots, 2n\}$. We create a positive e -requirement $\{x_1^{(e)}, \dots, x_n^{(e)}\}$ and a negative e -requirement $\{y_j^{n-1}\}$;

II.2.2. $\forall p_{1 \leq p \leq n-1}$ (ψ_p has a subterm-constant with value y_j^p , $y_j^p \notin \bigcup_{i \in N} A_i^{(s-1)}$

$\Rightarrow y_j^p \in \{x_1^{(e)}, \dots, x_n^{(e)}\}$). We find the terms $\psi^1, \dots, \psi^{n-1}$ corresponding to the terms $\tau^1, \dots, \tau^{n-1}$ according to Definition (3.6). From Theorem 1 for the terms $\tau^1, \dots, \tau^{n-1}$ there exist natural numbers i_1, \dots, i_n such that

$$i_1 \in M \ \& \ \dots \ \& \ i_n \in M \iff \tau_U^1(i_1, \dots, i_n) \in M \ \& \ \dots \ \& \ \tau_U^{n-1}(i_1, \dots, i_n) \in M$$

does not hold, i.e. one of the following holds:

$$(3.13) \ i_1 \in M \ \& \ \dots \ \& \ i_n \in M \ \& \ (\tau_U^1(i_1, \dots, i_n) \notin M \vee \dots \vee \tau_U^{n-1}(i_1, \dots, i_n) \notin M),$$

$$(3.14) \ (i_1 \notin M \vee \dots \vee i_n \notin M) \ \& \ \tau_U^1(i_1, \dots, i_n) \in M \ \& \ \dots \ \& \ \tau_U^{n-1}(i_1, \dots, i_n) \in M.$$

Let $m_j = \tau_U^j(i_1, \dots, i_n)$, $1 \leq j \leq n - 1$. We take: $A_j^{(s)} = A_j^{(s-1)} \cup \{x_j^{(e)}\}$, $j \in \{i_1, \dots, i_n\}$, and $A_l^{(s)} = A_l^{(s-1)}$ for $l \notin \{i_1, \dots, i_n\}$.

If (3.13) holds, then by Definition (3.3c) we have

$$x_1^{(e)} \in A^{(s)} \ \& \ \dots \ \& \ x_n^{(e)} \in A^{(s)}$$

and from Lemma (3.7):

$$z_i = \psi_V^i(x_1^{(e)}, \dots, x_n^{(e)}) \in [A_{m_j}^{(s)}],$$

and therefore $z_1 \notin A^{(s)} \vee \dots \vee z_{n-1} \notin A$, so in this case (i) is true.

If (3.14) holds, then $m_1 \in M \ \& \ \dots \ \& \ m_{n-1} \in M$ and therefore $z_1 \in A^{(s)} \ \& \ \dots \ \& \ z_{n-1} \in A^{(s)}$. We have also $i_1 \notin M \vee \dots \vee i_n \notin M$ and then $x_1^{(e)} \notin A^{(s)} \vee \dots \vee x_n^{(e)} \notin A^{(s)}$, i.e. (ii) is true.

We create a positive e -requirement $\{x_1^{(e)}, \dots, x_n^{(e)}\}$, which is also constant.

At the end we take $A_i = \bigcup_{s=0}^{\infty} A_i^{(s)}$, $i \in N$. Now we have to prove that the set

A satisfies the condition of our problem for n .

(3.15) Lemma. *For any e -number of a p.r. function the condition ((i) \vee (ii)) is injured only $2^e - 1$ times, i.e. we create not more than 2^e e -requirements.*

Proof. We shall use induction on e .

1) For $e = 0$ we have that the e -requirement can not be injured, because there is no requirement with higher priority and $2^0 - 1 = 0$.

For $e = 1$ we can injure the e -requirement only once when we create the single 0-requirement (if it exists) and $2^1 - 1 = 1$.

2) Let the statement hold for all the numbers smaller than e .

3) Let e be a number of a p.r. function. The condition ((i) \vee (ii)) for e is injured when a requirement with number between 0 and $e - 1$ is created, i.e. not more than $2^0 + 2^1 + \dots + 2^{e-1} = 2^e - 1$ times.

The lemma is proved.

(3.16) **Lemma.** *The set $N_1 \setminus A$ is infinite.*

Proof. Let $(N_1)_x = \{y \mid y \in N_1 \ \& \ y < x\}$. We shall prove that the set $(N_1)_{r(x)} \cap (N_1 \setminus A)$ contains at least x elements.

$$(N_1)_{r(x)} = \{y \mid y \in N_1 \ \& \ y < r(x)\} = \{y \mid y \in N_1 \ \& \ y < r'(n \cdot 2^x + x)\}$$

and because $r'(0) < r'(1) < \dots < r'(n \cdot 2^x + x - 1) < r'(n \cdot 2^x + x) < \dots$, the elements of the set $(N_1)_{r(x)}$ are $n \cdot 2^x + x$. Between them only 0-, 1-, \dots , $x - 1$ -requirements may be elements of A (the others are bigger than $r(x)$), i.e. not more than $n \cdot 2^0 + n \cdot 2^1 + \dots + n \cdot 2^{x-1} = 2^x$. Therefore, the elements of $(N_1)_{r(x)} \cap (N_1 \setminus A)$ are at least $n \cdot 2^x + x - n \cdot 2^x = x$ and the lemma is proved.

(3.17) **Lemma.** *For any natural number e such that $N_1 \subseteq \text{dom } \varphi_e$ (and especially for any e which is a number of recursive function) there exists a constant e -requirement.*

Proof. Let $e \in N$ and $N_1^n \subseteq \text{dom } \varphi_e$. Let us assume that there is not a constant e -requirement. We find s_0 such that at the step s_0 all e_1 -requirements for $e_1 < e$ are already built.

From Lemma (3.16) we have that there exist $x_1 \in N_1 \setminus A, \dots, x_n \in N_1 \setminus A$ such that $x_1 > r(e) \ \& \ \dots \ \& \ x_n > r(e)$. Let $s > s_0$ and $J^n(x_1, \dots, x_n) \in \text{dom } \varphi_{e,s}$. There at the step s a constant e -requirement is created.

The lemma is proved.

Let A be obtained from $\{A_i\}_{i \in N}$ according to Definition (3.3). According the construction, A is a r.e. set, $A^n \not\equiv_m A^{n-1}$, and $A^{n+1} \leq_m A^n$, i.e. the needed set is built.

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